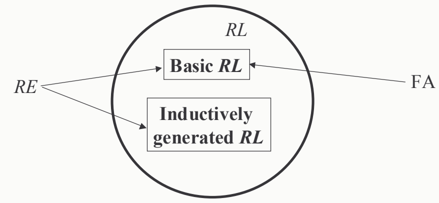
Session 7

Set Operations and FA

Descriptive power of *RE* and FA



We can define individual FA, so far!
But ee need to generate machines inductively!

Operations on FA

- According to Kleen's Theorem:
  - If L_1 and L_2 are both *RE* over Σ if and only if there are M_1 and M_2 accepting L_1 and L_2
- Definition of *RL*: IF L_1 and L_2 are *RL*
  - $L_1 \cup L_2$, L_1L_2 and L_1^* are also *RL*
- Set operations on *RL*
  - Union: $L_1 \cup L_2$ is a *RL*
  - Intersection: $L_1 \cap L_2$ is a *RL*
  - Difference: $L_1 - L_2$ is a *RL*
- Is there a way to generate machines for these languages from M_1 and M_2 (the machines accepting L_1 and L_2)?

Set operations on FA

- The case for $L_1 \cup L_2$
  - Consider whether a string $x \in L_1 \cup L_2$
  - It does if x belongs to either L_1 or L_2
  - Let M_1 and M_2
    $M_1 = (Q_1, \Sigma, q_1, A_1, \delta_1)$ such that $L(M_1) = L_1$
    $M_2 = (Q_2, \Sigma, q_2, A_2, \delta_2)$ such that $L(M_2) = L_2$
  - Process x with both M_1 and M_2
  - If x is accepted by either M_1 or M_2, it belongs to $L_1 \cup L_2$

Set operations on FA

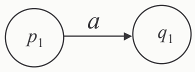
- Process the string x **simultaneously** by M_1 and M_2
  - $\delta_1^*(q_1, x) = p$
  - $\delta_2^*(q_2, x) = q$
- The three machines:
  - If $p \in A_1$ OR $q \in A_2$ then x is accepted by the union of M_1 and M_2
  - If $p \in A_1$ AND $q \in A_2$ then x is accepted by the intersection of M_1 and M_2
  - If $p \in A_1$ AND $q \notin A_2$ then x is accepted by the difference M_1 but not M_2

Set operations on FA

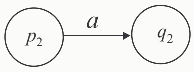
- The process step by step:
  - $\delta_1(p, a) = r$
  - $\delta_2(q, a) = s$
- The composite machine can be though in terms of abstract pairs of states (p, q) on the symbol a:
  - It moves from state (p, q) to $(\delta_1(p, a), \delta_2(q, a))$

Abstract states

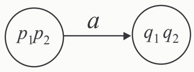
- Transition in FA-1:



- Transition in FA-2:



- Transition in comp. FA:



Theorem (constructive)

- Suppose
  - $M_1 = (Q_1, \Sigma, q_1, A_1, \delta_1)$ such that $\mathbf{L}(M_1) = L_1$
  - $M_2 = (Q_2, \Sigma, q_2, A_2, \delta_2)$ such that $\mathbf{L}(M_2) = L_2$
- Let $M = (Q, \Sigma, q_0, A, \delta)$ be a construction from M_1 and M_2 as follows:
  - $Q = Q_1 \times Q_2$
  - $q_0 = (q_1, q_2)$
  - For any $p \in Q_1$, $q \in Q_2$ and $a \in \Sigma$:

    $\delta((p, q), a) = (\delta_1(p, a), \delta_2(q, a))$
- Then
  - If $A = \{(p, q) \mid p \in A_1 \text{ OR } q \in A_2\}$ then M accepts $L_1 \cup L_2$
  - If $A = \{(p, q) \mid p \in A_1 \text{ AND } q \in A_2\}$ then M accepts $L_1 \cap L_2$
  - If $A = \{(p, q) \mid p \in A_1 \text{ AND } q \notin A_2\}$ then M accepts $L_1 - L_2$

Proof

- Acceptance of M_1 and M_2 is defined in terms of δ_1^* and δ_2^*, and acceptance by M is defined in terms of δ^*:
  - For $x \in \Sigma^*$ and $(p, q) \in Q$:

    $\delta^*((p, q), x) = (\delta_1^*(p, x), \delta_2^*(q, x))$
- A string x is accepted by M if and only if

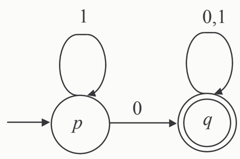
  $\delta^*((q_1, q_2), x) \in A$

  (i.e. Takes M from the initial to an accepting state)
- This is true if and only if

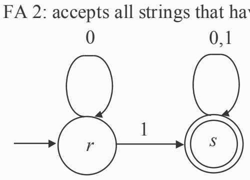
  $(\delta_1^*(q_1, x), \delta_2^*(q_2, x)) \in A$
- The set A is defined for case 1 as saying that
  - $\delta_1^*(q_1, x) \in \boldsymbol{A}_1$ OR $\delta_2^*(q_2, x) \in A_2$
  - Similarly for the other two cases

Example: Find the product FA

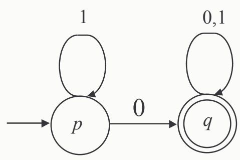
FA 1: accepts all strings that have a 0

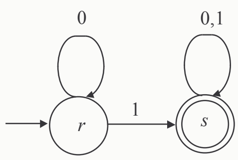


FA 2: accepts all strings that have a 1



Example: Find the product FA

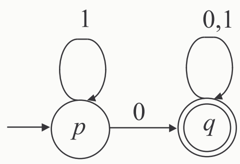


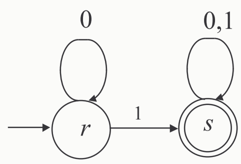


$Q = Q_1 \times Q_2$

s		
r		
	p	q

Example: Find the product FA

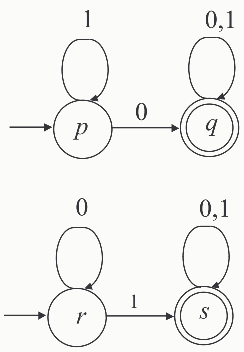




Transition table on 0

s		
r	qr	
	p	q

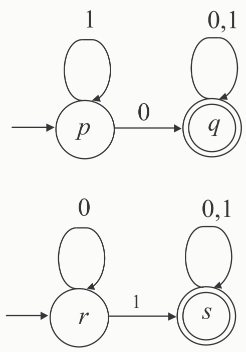
Example: Find the product FA



Transition table on 0

s	*qs*	
r	*qr*	
	p	*q*

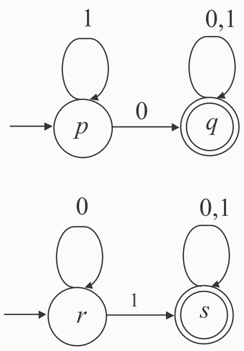
Example: Find the product FA



Transition table on 0

s	*qs*	
r	*qr*	*qr*
	p	*q*

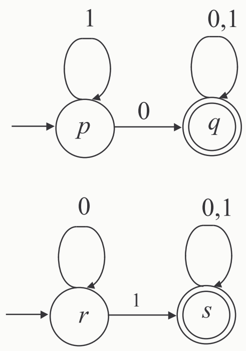
Example: Find the product FA



Transition table on 0

s	*qs*	*qs*
r	*qr*	*qr*
	p	*q*

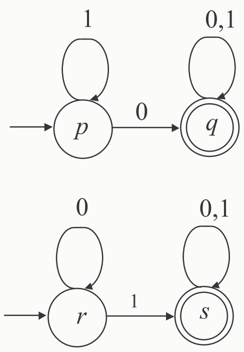
Example: Find the product FA



Transition table on 1

s	*ps*	*qs*
r	*ps*	*qs*
	p	*q*

Example: Find the product FA



Transition table on 0

s	*qs*	*qs*
r	*qr*	*qr*
	p	*q*

Transition table on 1

s	*ps*	*qs*
r	*ps*	*qs*
	p	*q*

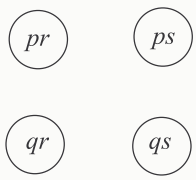
Theorem (constructive)

- Let M_1 and M_2 be
  - $M_1 = (Q_1, \Sigma, q_1, A_1, \delta_1)$ such that $L(M_1) = L_1$
  - $M_2 = (Q_2, \Sigma, q_2, A_2, \delta_2)$ such that $L(M_2) = L_2$
- Let $M = (Q, \Sigma, q_0, A, \delta)$ be a construction from M_1 and M_2 as follows:
  - The set of states: $Q = Q_1 \times Q_2$
  - The initial state: $q_0 = (q_1, q_2)$
  - For any $p \in Q_1$, $q \in Q_2$ and $a \in \Sigma$

    $\delta((p, q), a) = (\delta_1(p, a), \delta_2(q, a))$
- So, far we have achieved this!

Find the graph of FA

The set of states:



Transition table on 0

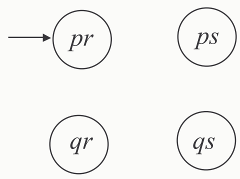
s	*qs*	*qs*
r	*qr*	*qr*
	p	*q*

Transition table on 1

s	*ps*	*qs*
r	*ps*	*qs*
	p	*q*

Find the graph of FA

The initial state:



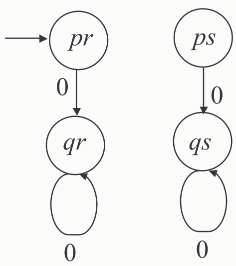
Transition table on 0

s	*qs*	*qs*
r	*qr*	*qr*
	p	*q*

Transition table on 1

s	*ps*	*qs*
r	*ps*	*qs*
	p	*q*

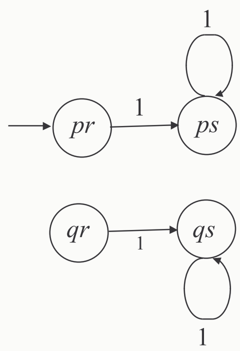
The transition function of FA



Transition table on 0

s	*qs*	*qs*
r	*qr*	*qr*
	p	*q*

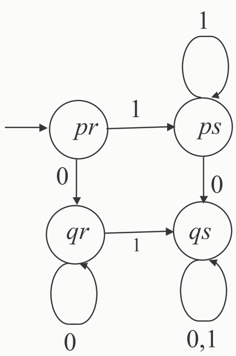
The transition function of FA



Transition table on 1

s	*ps*	*qs*
r	*ps*	*qs*
	p	*q*

The transition function of FA



Transition table on 0

s	*qs*	*qs*
r	*qr*	*qr*
	p	*q*

Transition table on 1

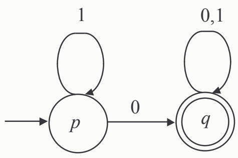
s	*ps*	*qs*
r	*ps*	*qs*
	p	*q*

The accepting states

- The final states:
  - If $A = \{(p, q) \mid p \in A_1 \text{ OR } q \in A_2\}$ then M accepts $L_1 \cup L_2$
  - If $A = \{(p, q) \mid p \in A_1 \text{ AND } q \in A_2\}$ then M accepts $L_1 \cap L_2$
  - If $A = \{(p, q) \mid p \in A_1 \text{ AND } q \notin A_2\}$ then M accepts $L_1 - L_2$

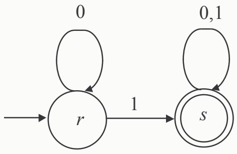
Accepting states FA

FA 1: accepts all strings that have a 0



$A_1 = \{q\}$

FA 2: accepts all strings that have a 1



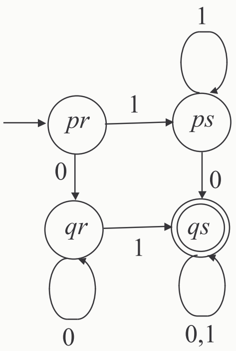
$A_2 = \{s\}$

Find the Intersection

- If A = $\{(p, q) \mid p \in A_1$ AND $q \in A_2\}$
  then M accepts $L_1 \cap L_2$
  - $Q = \{pr, ps, qr, qs\}$
  - $A_1 = \{q\}$ AND $A_2 = \{s\}$

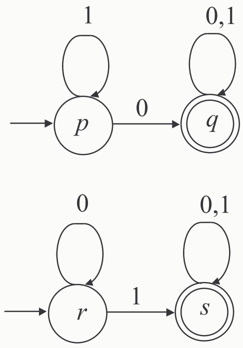
  then
  - $A = \{qs\}$

The Intersection FA



Accepts all strings including at least one 0 and one 1

But suppose we want the union FA!



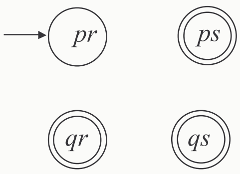
The language consisting of at least one 0 or at least one 1

Find the Union FA

- If A = $\{(p, q) \mid p \in A_1$ OR $q \in A_2\}$
  then M accepts $L_1 \cup L_2$
  - $Q = \{pr, ps, qr, qs\}$
  - $A_1 = \{q\}$ OR $A_2 = \{s\}$

  then
  - $A = \{ps, qr, qs\}$

Find the Union FA



Final states have *q* or *s*!

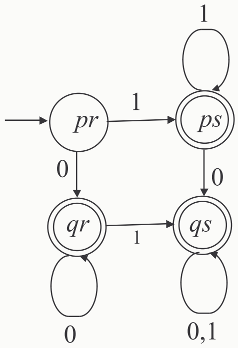
Transition table on 0

s	*qs*	*qs*
r	*qr*	*qr*
	p	*q*

Transition table on 1

s	*ps*	*qs*
r	*ps*	*qs*
	p	*q*

The Union FA



Accepts all strings including at least one 0, or at least one 1 and at least one 0 and 1 one!

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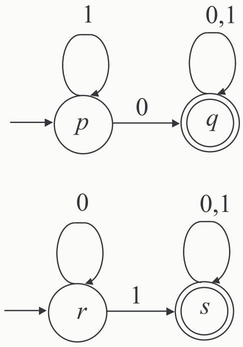
Or the difference FA!



The language consisting of strings containing at least one 0 but NOT at least one 1

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Find the Difference FA

- If $A = \{(p, q) \mid p \in A_1$ AND $q \notin A_2\}$ then M accepts $L_1 - L_2$
  - $Q = \{pr, ps, qr, qs\}$
  - $A_1 = \{q\}$ BUT NOT $A_2 = \{s\}$

  then
  - $A = \{qr\}$

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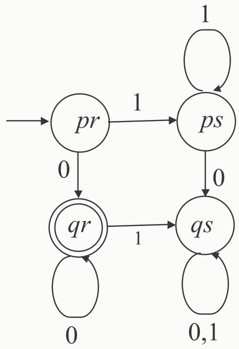
The Difference FA



Accepts all strings including at least one 0, but not strings including one 1!

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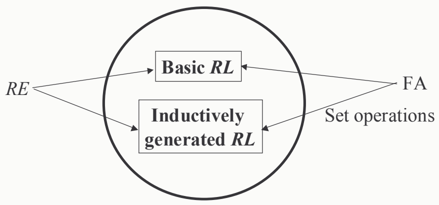
Descriptive power of *RE* and FA



But we still need concatenation and closure!

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